#### 1 Coding (Favorite Monte Carlo algorithm).

Since this class is wrapping up, what is your favorite Monte Carlo algorithm and why? Provide a numerical illustration.

### 2 Coding (Randomized subspace iteration).

Apply the randomized SVD algorithm (randomized subspace iteration with q = 1) to approximate the diagonal matrix

$$\mathbf{A} = \text{diag}(e^{-.1}, e^{-.2}, \dots, e^{-999.9}).$$

The approximation will be quite accurate if you use a block size  $k \ge 50$ . Next, try to approximate a matrix **B** which is a small entrywise perturbation of **A**:

$$b_{ij} = a_{ij} + Z_{ij}, \qquad Z_{ij} \stackrel{\text{ind}}{\sim} \mathcal{N}(0, 0.002^2).$$

Approximating B will be hard, necessitating a large block size k and/or number of iterations q. Why?

## 3 Coding (Randomly pivoted Cholesky)

Randomly pivoted Cholesky selects "diverse" columns to approximate a matrix  $\boldsymbol{A}$ . To see this, generate random data points  $\boldsymbol{x}^{(1)}, \ldots, \boldsymbol{x}^{(N)}$  from a mixture of Gaussians with centers at (+2, +2), (+2, -2), (-2, +2), and (-2, -2) and variance 1/4. Define the kernel matrix  $\boldsymbol{A} \in \mathbb{R}^{N \times N}$  with entries  $a_{ij} = \exp(-\|\boldsymbol{x}^{(i)} - \boldsymbol{x}^{(j)}\|^2/2)$ , which quantifies the similarity between data points (close to one means similar, close to zero means dissimilar). Apply randomly pivoted Cholesky to approximate the matrix  $\boldsymbol{A}$ . Which data points correspond to the first 4 selected rows? How similar is the approximate matrix  $\hat{\boldsymbol{A}}$  to the target matrix  $\boldsymbol{A}$ ?

# 4 Coding (k-means++).

If you have ever used k-means to cluster the rows of a data matrix  $\mathbf{A} \in \mathbb{R}^{L \times N}$ , you have used randomized numerical linear algebra. The standard initialization for k-means, called k-means++ builds up a randomized rank-k approximation  $\hat{\mathbf{A}} \in \mathbb{R}^{L \times N}$ . We initialize the approximation by setting  $\hat{\mathbf{A}} = \mathbf{0}$ . Then, we perform k updates by randomly selecting a row index  $s \in \{1, \ldots, L\}$  with

$$\mathbb{P}\{s=i\} = \sum_{j=1}^{N} |a_{ij} - \hat{a}_{ij}|^2.$$

and update each row of  $\hat{A}$  as  $\hat{a}_{ij} = a_{sj}$  for  $1 \leq j \leq N$  if

$$\sum_{j=1}^{N} |a_{ij} - a_{sj}|^2 \le \sum_{j=1}^{N} |a_{ij} - \hat{a}_{ij}|^2.$$

Apply k-means++ to the kernel matrix  $\mathbf{A} \in \mathbb{R}^{N \times N}$  constructed in problem #3. Which data points correspond to the first 4 selected rows? How similar is the approximate matrix  $\hat{\mathbf{A}}$  to the target matrix  $\mathbf{A}$ ? Compare and contrast with randomly pivoted Cholesky.

## 5 Choose-your-own-adventure (Monte Carlo connection).

Can we view randomized subspace iteration and randomly pivoted Cholesky as Monte Carlo algorithms? If so, how?

## 6 Math (Nyström approximation for psd matrices).

To approximate a general, rectangular matrix  $A \in \mathbb{R}^{L \times N}$ , we can form a low-rank approximation

$$\hat{A} = \Pi_X A = XY^*, \qquad Y = A^*X, \tag{1}$$

where  $X \in \mathbb{R}^{L \times k}$  is a rank-k orthogonal matrix. If  $A \in \mathbb{R}^{N \times N}$  is positive semidefinite (psd), we can alternatively form the Nyström approximation

$$\hat{\boldsymbol{A}} = \boldsymbol{A}^{1/2} \boldsymbol{\Pi}_{\boldsymbol{A}^{1/2} \boldsymbol{X}} \boldsymbol{A}^{1/2} = \boldsymbol{Y} (\boldsymbol{X}^* \boldsymbol{Y})^{\dagger} \boldsymbol{Y}^*.$$
<sup>(2)</sup>

Which approximation leads to a smaller Frobenius norm error, (1) or (2)? Provide a proof.

### 7 Math (Largest eigenvalue).

Let's prove some bounds for randomized subspace iteration.

(a) If we use randomized subspace iteration to approximate the largest singular value  $\sigma_1(\mathbf{A})$  of a matrix  $\mathbf{A} \in \mathbb{R}^{L \times N}$ , prove that

$$\hat{\sigma}_1(\boldsymbol{A})^2 = \max_{\boldsymbol{\omega} \in \operatorname{range}(\boldsymbol{\Omega})} \frac{\boldsymbol{\omega}^* (\boldsymbol{A}^* \boldsymbol{A})^{2q} \boldsymbol{\omega}}{\boldsymbol{\omega}^* (\boldsymbol{A}^* \boldsymbol{A})^{2q-1} \boldsymbol{\omega}},\tag{3}$$

where  $\mathbf{\Omega} \in \mathbb{R}^{N \times k}$  is the Gaussian initialization matrix.

- (b) Show that the distribution of  $\hat{\sigma}_1(\mathbf{A})^2$  only depends on the singular values of  $\mathbf{A}$ , regardless of the singular vectors.
- (c) Now assume without loss of generality that A is diagonal and psd, with non-increasing diagonal entries and rewrite (3) in a simpler form.
- (d) Partition  $\Omega = \begin{bmatrix} \Omega_1 \\ \Omega_2 \end{bmatrix}$  and set  $\omega$  to be the first column of  $\Omega \Omega_1^{\dagger}$ . Use (3) to obtain write down relatively simple, explicit lower and upper bounds on  $\hat{\sigma}_1(\mathbf{A})^2$ .
- (e) What random matrix theory do we need to bound the expected value of  $\hat{\sigma}_1(\mathbf{A})^2$ ?